

**IT**

**SECURITY**

RESEARCH GROUP

## **Secure and Efficient Handover Protocols for WMNs**

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Preliminaries

Related Mechanisms

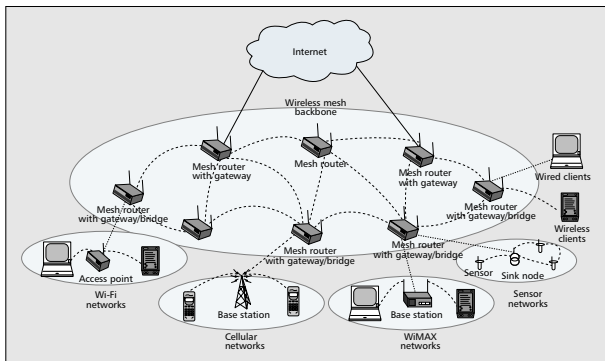
Approach

Security Considerations

Conclusion

# Motivation

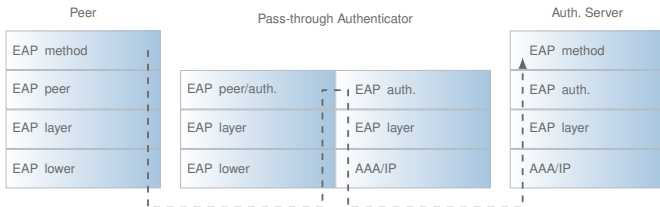
Enabling handover with infrastructure support in a secure fashion



- ▶ WMNs are easy to deploy, to extend, and provide relatively high data rates
- ▶ Existing security architectures seldom address secure handover
- ▶ How can we leverage the infrastructure of these networks?

# EAP

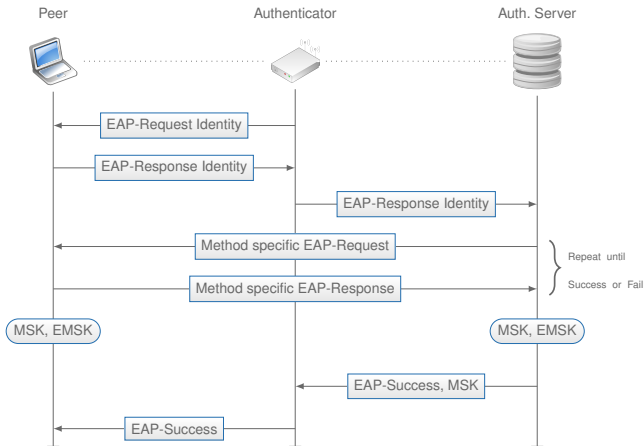
## Extensible Authentication Protocol



- ▶ Supports multiple authentication methods (EAP-TTLS, PEAP, AKA,...)
- ▶ Key generating EAP methods:
  - ▶ MSK (Master Session Key)
  - ▶ EMSK (Extended Master Session Key)

## EAP

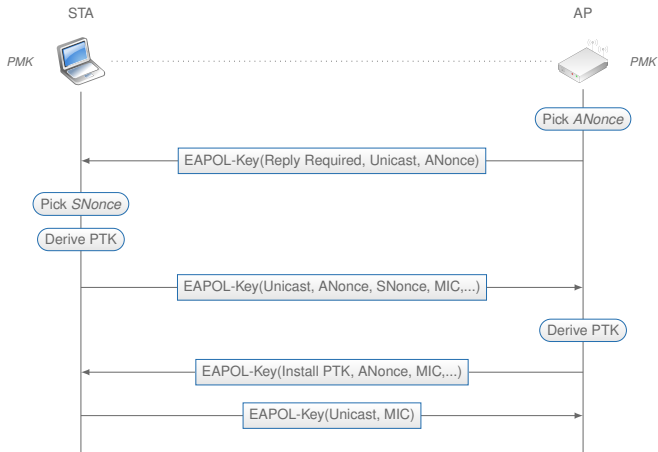
## EAP Operation



- ▶ MSK and EMSK are generated by EAP method
- ▶ PMK = first 256 bits of MSK
- ▶ 4-way handshake uses PMK to derive transient keys

## IEEE 802.11i

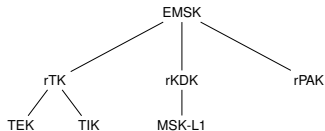
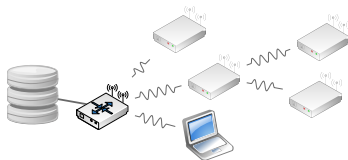
## 4-way handshake



- ▶ Goal: Secure link layer communication
- ▶ EAPOL = EAP Transport Over LAN
- ▶ ~ 4ms

# FSASD

## Framework for Sequential Deployment of WMNs



- ▶ Allows simple and secure sequential deployment of WMNs
- ▶ Uses EAP-TTLS which provides secure EAP authentication over several hops
- ▶ EMSK spawns key hierarchy to derive keys for IPsec ESP
- ▶ Authentication traffic is secured by IPsec
- ▶ 3-party handshake protocol for sequential deployment (3PHSD)
  - ▶ MC ↔ MG, MC ↔ MC, MC ↔ MR

## Related Mechanisms

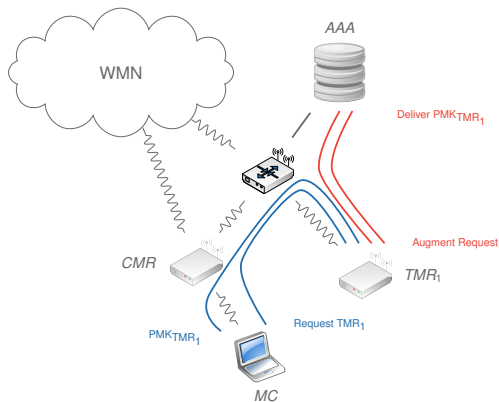
...and "standards" *not* for WMNs

- ▶ 802.11r - trusted APs, no secure key transport between APs
- ▶ ERP - insecure key transport AAA-NAS
- ▶ CAPWAP - new split of PHY/MAC, custom devices required



## 3-Party-Handshake Protocol for Handover (3PHSH)

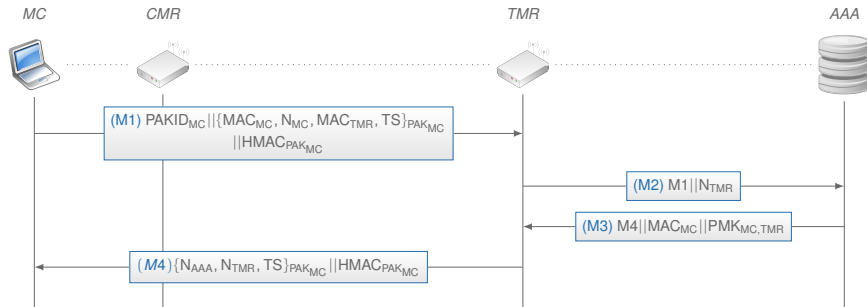
### Overview



- ▶ Extend 3PHSD to pro-actively request a fresh PMK for one TMR
- ▶ TMR and MC cache PMK
- ▶ Use cached PMK to avoid full EAP-TTLS re-authentication

# 3-Party-Handshake Protocol for Handover (3PHSH)

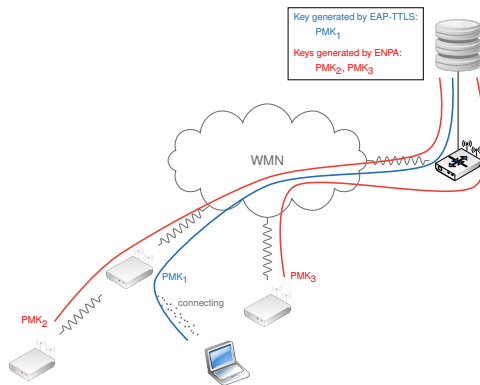
## Message Flow



$$\text{PMK}_{MC,TMR} = \text{KDF}(r\text{KDK}_{MC,AAA}, \text{key\_label} || N_{TMR} || N_{AAA} || N_{MC} || \text{MAC}_{MC} || \text{MAC}_{TMR})$$

# EAP-TTLS Neighborhood Pre-Authentication (ENPA)

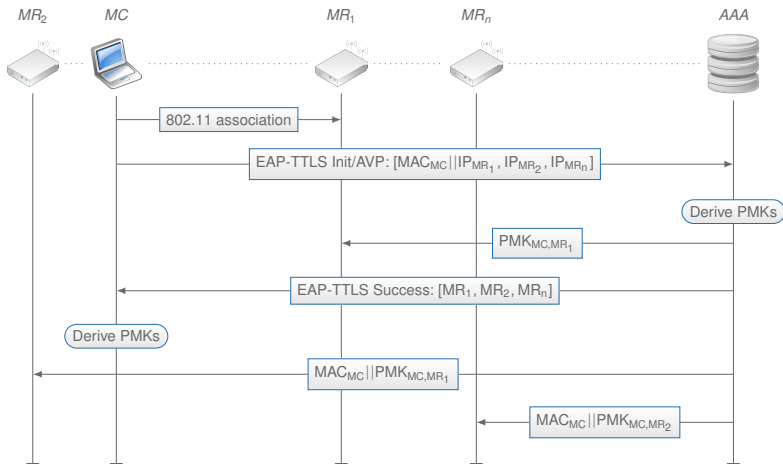
## Overview



- ▶ Request fresh PMKs for the other MRs during initial association
- ▶ Embedded into EAP-TTLS authentication
- ▶ Also works with EAP methods supporting Diameter AVPs
- ▶ Use cached PMKs to avoid full EAP-TTLS re-authentication

## EAP-TTLS Neighborhood Pre-Authentication (ENPA)

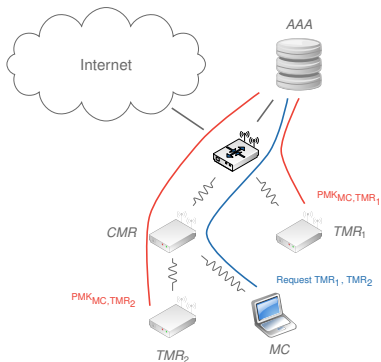
## Message Flow



$$PMK_{MC, TMR_i} = KDF(rKDK_{MC, AAA}, key\_label || IP_{TMR_i} || tls\_client\_random || tls\_server\_random)$$

# Neighborhood Pre-Authentication (NPA)

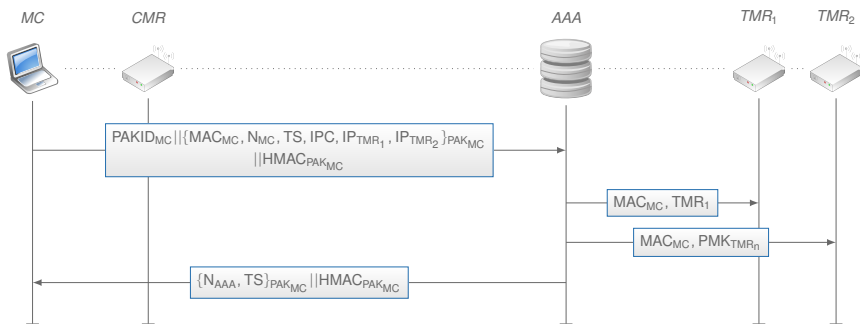
## Overview



- ▶ “Combine” 3PHSH and ENPA
- ▶ MC is connected to WMN
- ▶ Request PMKs for several TMRs in range at once
- ▶ Use cached PMK to avoid full EAP-TTLS re-authentication

# Neighborhood Pre-Authentication (NPA)

## Message Flow



$$\text{PMK}_{MC, TMR_j} = \text{KDF}(r\text{KDK}_{MC, AAA}, \text{key\_label} || \text{MAC}_{MC} || N_{MC} || N_{AAA} || \text{IP}_{TMR_j}) \quad (1)$$

# Security Analysis

## Overview

- ▶ Analysis is based on “Best Practice” RFC 4962 “Guidance for AAA Key Management”
- ▶ 11 requirements a protocol should satisfy
- ▶ ENPA is embedded in a AAA protocol
- ▶ 3PHSH and NPA can only be used after a successful authentication

## Security Analysis

1. Cryptographic algorithm independent ✓
2. Strong, fresh session keys ✓
3. Limit key scope ✓
4. Replay Protection ✓
5. Authenticate all parties ✓
6. Peer and authenticator authorization ✓
7. Keying material confidentiality and integrity ✓
8. Confirm ciphersuite selection ( )
9. Uniquely named keys ✓
10. Prevent the domino effect ✓
11. Bind the key to its context ✓



## Performance Evaluation

See paper

## Conclusion

- ▶ No secure and practical approaches for WMNs exist
- ▶ Three protocols to pro-actively request PMKs
  - ▶ 3PHSH
  - ▶ ENPA
  - ▶ NPA
- ▶ MC determines candidate MRs itself (currently IP included in beacons)
- ▶ Exploited the security properties of FSASD
- ▶ Two protocols are faster than the EAP-TTLS re-authentication
  - ▶ 3PHSH (wireless hops 7):  $\approx 25ms + 4ms$
  - ▶ NPA (max. distance 4, IP count 5):  $\approx 10ms + 4ms$
- ▶ ENPA nearly same duration as EAP-TTLS:  $300ms$ 
  - ▶ takes  $5ms$  longer for adding 1 TMR
  - ▶ only 4-way-handshake needed to roam afterwards:  $4ms$

Q?

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